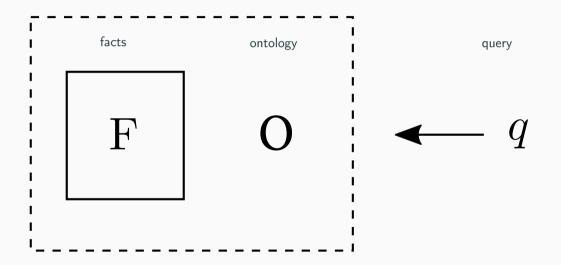
Représentation des connaissances et raisonnement

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06/04/2023 - Journée LIMOS

Ontology-based Query Answering Problem

knowledge base



Facts	Ontology	Query
Human(Alice)	$\mathit{Human}(x) \land \mathit{ParentOf}(y,x) \rightarrow \mathit{Human}(y)$	Human(Bob) ?
ParentOf(Bob, Alice)	$\mathit{Human}(x) o \exists z \; \mathit{ParentOf}(z,x)$	

Facts	Ontology	Query
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Facts	Ontology	Query
Human(Alice)	$\mathit{Human}(x) \land \mathit{ParentOf}(y,x) ightarrow \mathit{Human}(y)$	$\exists u \; ParentOf(u, Bob) ?$
ParentOf (Bob, Alice)	$\mathit{Human}(x) o \exists z \; \mathit{ParentOf}(z,x)$	

Ontology	Query
$\mathit{Human}(x) \land \mathit{ParentOf}(y,x) \rightarrow \mathit{Human}(y)$	$\exists u \; ParentOf(u,u) ?$
$\mathit{Human}(x) o \exists z \; \mathit{ParentOf}(z,x)$	
	$Human(x) \wedge ParentOf(y,x) \rightarrow Human(y)$

Conjunctive Query Answering using Existential Rules

An existential rule is a FOL formula:

$$\forall \bar{x} \ \forall \bar{y} \ B(\bar{x}, \bar{y}) \rightarrow \exists \bar{z} \ H(\bar{x}, \bar{z})$$

where B and H are conjunctions of atoms

An undecidable problem

- F a set of relationnal facts
- \bullet \mathcal{R} a set of existential rules
- q a conjunctive query

 F, \mathcal{R} entails q?

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- F a set of relationnal facts
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 F, \mathcal{R} entails q?

What are the class of rule sets for which QA is always decidable?

Saturation of facts

Facts	Rules	Query
Human(Alice)	$\mathit{Human}(x) \land \mathit{ParentOf}(y,x) ightarrow \mathit{Human}(y)$	Human(Bob) ?
ParentOf (Bob, Alice)	$\mathit{Human}(x) o \exists z \; \mathit{ParentOf}(z,x)$	
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Human(Bob)		
$ParentOf(n_0, Bob)$		

Saturation of facts

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Human(Bob)		
$ParentOf(n_0, Bob)$		
$ParentOf(n_1, n_0)$		
•••		

Saturation of facts

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We can decide the query answering problem if the saturation is finite.

Facts	Rules	Query
Human(Alice)	$\mathit{Human}(x) \land \mathit{ParentOf}(y,x) \rightarrow \mathit{Human}(y)$	Human(Bob)
ParentOf (Bob, Alice)		?

Facts	Rules	Query
Human(Alice)	$\mathit{Human}(x) \land \mathit{ParentOf}(y,x) \rightarrow \mathit{Human}(y)$	Human(Bob)
ParentOf(Bob, Alice)		$\vee \exists x \; Human(x) \wedge ParentOf(Bob, x)$
		?

Facts	Rules	Query	
Human(Alice)	$AncestorOf(x,y) \land ParentOf(y,z) \rightarrow$	$\exists u \; AncestorOf(u, Bob)$?
ParentOf (Bob, Alice)	AncestorOf(x, z)		

Facts	Rules	Query
Human(Alice)	$AncestorOf(x,y) \land ParentOf(y,z) \rightarrow$	$\exists u \ AncestorOf(u, Bob)$
ParentOf (Bob, Alice)	AncestorOf(x, z)	$\forall \exists u, y \ AncestorOf(u, y) \land ParentOf(y, Bob)$?

Rewriting of the query

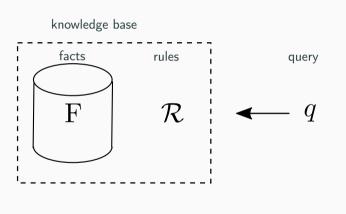
Facts	Rules	Query
Human(Alice)	$AncestorOf(x,y) \land ParentOf(y,z) \rightarrow AncestorOf(x,z)$	$\exists u \; Ancestor Of(u, Bob)$
${\it ParentOf}({\it Bob},{\it Alice})$		$\forall \exists u, y \ AncestorOf(u, y) \land ParentOf(y, Bob)$
		$\vee \exists u, y, y' \ \textit{AncestorOf}(u, y') \land$
		$ParentOf(y',y) \land$
		ParentOf(y, Bob)?

We can decide the query answering problem if the query rewriting is finite.

Rule Sets with Decidable Query Answering

- Finite saturation rule set : Datalog
- Finite rewriting rule set : DL_{lite}
- Bounded tree width rule set : most of the DLs

In Practice: Ontology-based Data Management



Materialization-based approach

- Time and stockage for saturation
- Good query time
- Updates require a saturation maintenance

Query rewriting approach

- Query rewriting may be large
- Updates for free

Current Research Topics

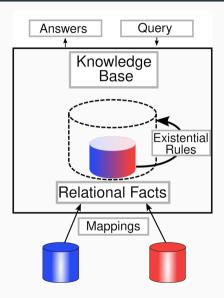
Handling infinite saturation

- Take advantage of regularity of saturation
- Create a finite representation of the infinite saturation

Pruning infinite rewriting

- Take advantage of the facts at hand
- Avoid unnecessary rewriting steps

Ontology-Based Data Access



- Integrating the facts coming from heterogeneous sources through mappings
- Use ontologies to manage the semantic heterogeneity
- Mediation-based approach :
 - query rewriting with the rules and mappings
 - 2. evaluation using mediator engine